

14th September 2025 in Kuala Lumpur, Malaysia

Improving Fault Vulnerability Detection via Rehosting and Comparative Analysis of Open-source Tools

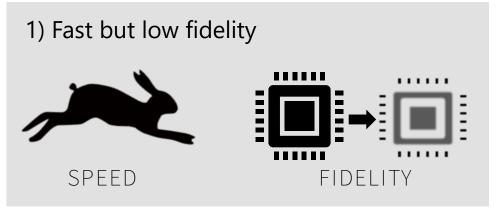
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This presentation is based on results obtained from a project, JPNP24003, commissioned by the New Energy and Industrial Technology Development Organization (NEDO).

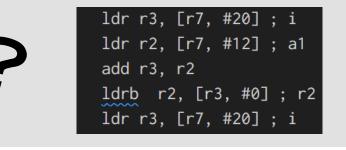
This Talk in Brief



Challenges





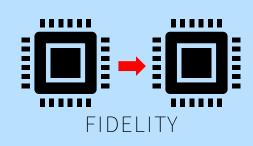


Solution

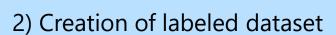


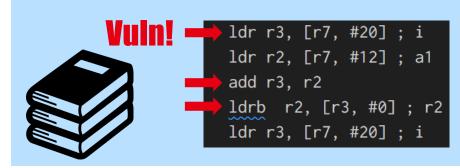












Results: Rehosting achieves near-perfect recall (99%-100%) while keeping high speed

Outline



- Background
- Fault Vulnerability Detection
- Rehosting
- Experiment
- Discussion
- Conclusion



Background



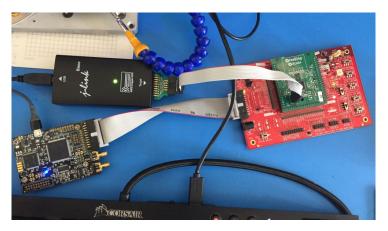
Fault Injection Attack (FIA) against Actual Product



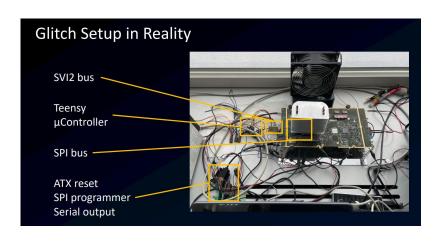
If software is secure, attack hardware



Game console [5]



Cryptocurrency wallet [6]



Automotive unit [4]

- Software security for embedded devices is improving
- Attackers seek the next "weakest point," targeting fault attacks

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^[4] C. Werling, N. K"uhnapfel, H. N. Jacob, and O. Drokin, "Jailbreaking an Electric Vehicle in 2023 or What It Means to Hotwire Tesla's x86-Based Seat Heater," BlackHat Asia, 2023.

^[5] plutoo, derrek and naehrwert, "Console Security – Switch", 2017

^[6] J. Datko, C. Quartier, and K. Belyayev, "Breaking Bitcoin Hardware Wallets: Glitches cause stitches!" DEF CON 2017, 2017

FIA Resistance in Software Design Phase



Increasing demand for fault vulnerability detection (FVD) from binary

Why software-based detection?

Scalable, reproducible, cost-effective comparing to physical FIA

Why binary?

- Higher-level languages diverge from physical mechanisms
- Binary is the closest to hardware within the software layer

```
8000054: 697b
                 ldr r3, [r7, #20]; i
8000056: 68fa
                 ldr r2, [r7, #12]; a1
8000058: 4413
                 add r3, r2
800005a: 781a
                 ldrb r2, [r3, #0]; r2 =
800005c: 697b
                 ldr r3, [r7, #20]; i
800005e: 68b9
                 ldr r1, [r7, #8]
8000060: 440b
                 add r3, r1
8000062: 781b
                 ldrb r3, [r3, #0]; r3 =
8000064: 429a
                 cmp r2, r3
8000066: d001
                 beq.n 800006c <byteArrayCo
8000068: 2300
                movs r3, #0
800006a: e007
                 b.n 800007c <byteArrayComp
```

Concept of FVD from binary

Challenges



1) Speed vs. Fidelity, 2) Lack of benchmark dataset

1. Trade-offs between speed and fidelity:

- Instruction Set Emulation (ISE)-based tools prioritize speed by omitting system-level emulation (→ QEMU)
- This simplification leads to inaccurate memory and register initialization, causing false positives and negatives

2. Lack of Standardized Benchmarks:

- Lack of ground-truth labeled datasets makes fair and reproducible tool comparison difficult
- This is due to the difficulty in manual verification of "all fault patterns"



2) Lack of benchmark dataset



```
ldr r3, [r7, #20] ; i
ldr r2, [r7, #12] ; a1
add r3, r2
ldrb r2, [r3, #0] ; r2
ldr r3, [r7, #20] ; i
```

Our Contributions



1) Rehosting for FVD, 2) Integration of results from multiple tools

1. Rehosting for Improved Fidelity:

- Introduced a rehosting technique to mitigate misclassification in ISE-based tools (e.g., FaultFinder).
- Accurately reproduces memory and register states, ensuring faithful system initialization

1) Rehosting improves fidelity SPEED FIDELITY

2. Labeled Benchmark Creation and Comparative Analysis

- Constructed a labeled benchmark dataset by manually verifying vulnerabilities from multiple tools
- This significantly reduces verification cost and enables systematic misclassification pattern identification
- This benchmark enables to evaluate the speed and detection performance of FVD tools

```
2) Creation of labeled

ldr r3, [r7, #20]; i
ldr r2, [r7, #12]; a1
add r3, r2
ldrb r2, [r3, #0]; r2
ldr r3, [r7, #20]; i
```



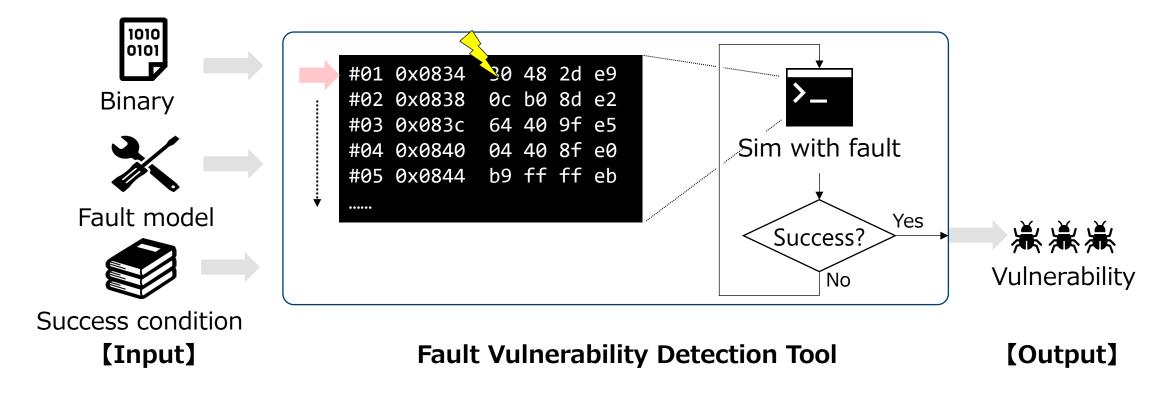
Fault Vulnerability Detection



Processing Flow of Fault Vulnerability Detection (FVD)



Run binary simulating fault and check if it satisfies success condition



- Vulnerabilities can be comprehensively detected by injecting faults into all instructions and verifying success conditions
- Example) Password authentication, DFA on AES



Characteristics of FVD Tools



Differences in tools make common evaluation difficult

Tool	Architecture	Platform	Success condition		
FaultFinder	ARM / x86 / RISC-V	Unicorn (ISE)	Target addr., reg/mem		
FaultArm	ARM / x86	- (Parser)	Anti-pattern		
Fault-injection-simulation*	ARM	angr	Target addr.		
ARCHIE	ARM / / RISC-V	QEMU -system	Target addr., reg/mem		
ARMORY	ARM	M-ulator (ISE)	Target addr.		
Chaos Duck	ARM / x86	Native/ QEMU -user	stdout / exit code		
FaultInjectionSimulator*	/ x86	Native	exit code		

*Not published as a paper

- Seven OSS-based FVD tools selected from GitHub
- Differences:
 - Architecture: Even if the architecture is the same, the user binary/system binary differs
 - Platform: ISE, QEMU, and native execution are primary
 - Success condition: target address, register/memory condition, exit code, ...



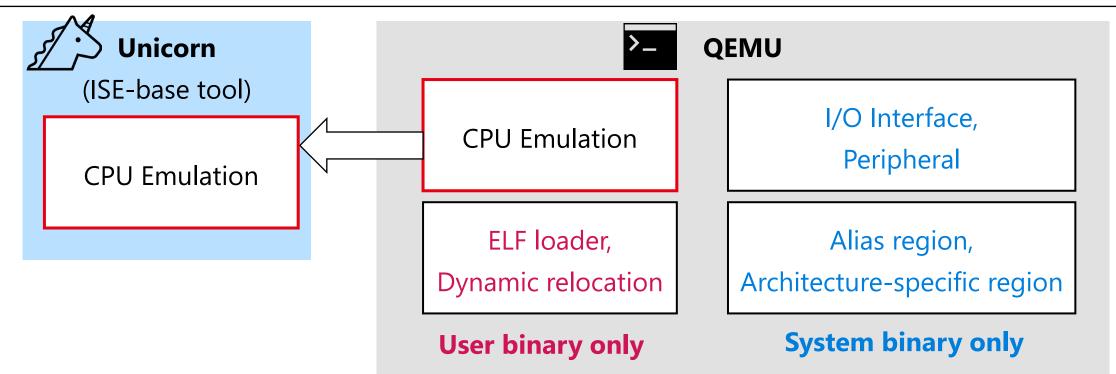
Rehosting



Problem: ISE-based Tools Lack Memory Reproducibility





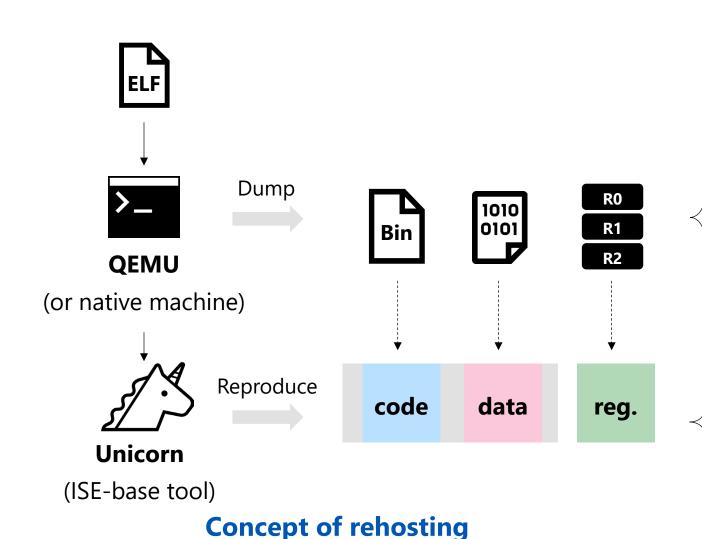


- ISE-based tools achieves high speed by emulating only execution of instruction set
- Omitted functions cause **differences in memory state** at entry point (*main()*)
 - → Leads to false positives and false negatives in vulnerability detection

Proposed: Rehosting



Porting the processor state (memory) during QEMU/native execution



1. Execute ELF on QEMU

2. Debug ELF with gdb

3. Break at entry point (main())

4. Dump memory (code and memory)

5. Dump CPU state (register)

6. Reproduce CPU state

Typical rehosting flow



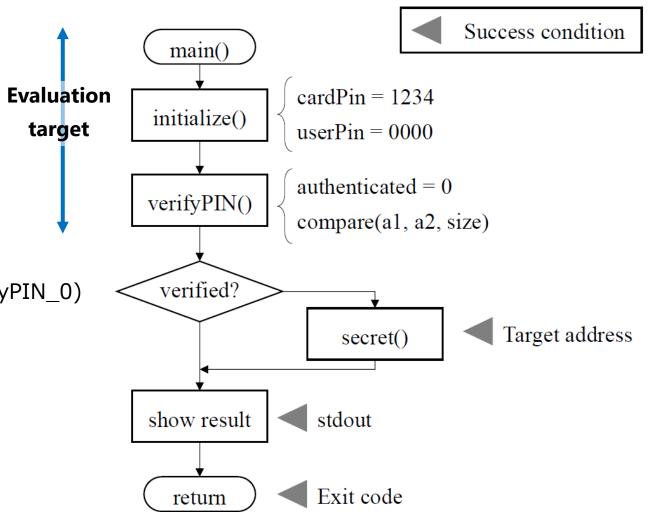
Experiment

Setup: Cross-Tool Evaluability Design



Case study: "PIN verification" as a common benchmark

- Target: ARMv7m, ARMv7a, x86
- Fault models:
 - Instruction skipping (IS)
 - Bit-flip on Instruction (BF-I)
 - Bit-flip on Register (BF-R)
- Benchmark program
 - PIN verification without countermeasure (VerifyPIN_0)
 from FISSC [12]
 - Hardcoded userPIN fails in PIN verification
 - Program modification:
 Support all success conditions
 (stdout / reached path / exit code)



Evaluation Metrics



Common vulnerability detection approach: maximizing recall while maintaining high precision

Metrics 1: Missed Faults

- Proportion of detected cases out of all vulnerabilities
- $Recall = \frac{TP}{TP + FN}$

Metrics 2: Effective Faults

- Proportion of actual vulnerabilities out of all detected cases
- $Precision = \frac{TP}{TP+FP}$

Ground Truth

 Integrating vulnerabilities reported by each tool, followed by manual analysis

Actual

	TRUE	FALSE
TRUE	TP True Positive	FP False Positive
FALSE	FN False Negative	TN True Negative

Result: Model1 Instruction Skipping



FaultFinder run on multi architectures with high performance and fast speed

- Rehosting eliminates all FNs, increasing recall up to 100%, while decreasing precision in some cases
- Rehosting keeps fast speed (low overhead)
- ISE-based tools are >50x faster than QEMU-based tools

		Tool*	Time [s]	Recall	Precision
		FaultFinder (rehost)	1.7	1.00	1.00
ARM	FaultFinder	1.6	0.88	0.94	
	v7a	FaultArm	2.0	0.59	0.29
		Chaos Duck	107.0	0.88	1.00
		FaultFinder (rehost)	1.7	1.00	0.89
		FaultFinder	1.6	0.94	1.00
ARM v7m	FaultArm	1.9	0.31	0.31	
	ARCHIE	173.7	1.00	0.89	
	ARMORY	0.4	0.75	1.00	
		FaultFinder (rehost)	1.4	1.00	1.00
		FaultFinder	1.3	1.00	1.00
x86	FaultArm	1.8	0.07	0.20	
	Chaos Duck	12.4	0.93	1.00	
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Both time and recall are good

Good

ISE

Bad

Others

* Only tools proposed in papers

QEMU

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Result: Model2,3 Bit-Flip on Instruction/Register



(Same trend) FaultFinder run on multi architectures with high performance and fast speed

(repeated)

- Rehosting eliminates all FNs, increasing recall up to 100%, while decreasing precision in some cases
- Rehosting keeps fast speed (low overhead)
- ISE-based tools are >400x faster than QEMU-based tools

ISE

	Tool*	Time [s]	Recall	Precision
ARM	FaultFinder (rehost)	8.9	0.99	0.95
v7a	FaultFinder	5.2	0.80	0.89
BF-I	Chaos Duck	3883.9	0.85	0.96
	FaultFinder (rehost)	4.4	1.00	0.85
ARM	FaultFinder	2.8	0.92	1.00
v7m BF-I	ARCHIE	8661.5	1.00	0.85
BF-I	ARMORY	0.4	0.70	0.93
	FaultFinder (rehost)	7.4	1.00	0.99
ARM –	FaultFinder	4.3	0.97	1.00
v7m BF-R	ARCHIE	8401.7	1.00	0.90
Dr-K	ARMORY	0.4	0.62	1.00

Both time and recall are good

Good

Bad

Others

QEMU

^{*} Only tools supporting bit-flip



Misclassification Analysis & Limitation



Misclassification Analysis



Misclassification fall into 10 patterns

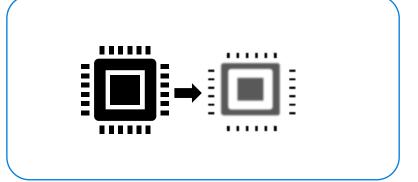
Tool	FP	FN	Init. state	Invalid access	Write attribute	Unimplemented inst.	Others
FaultFinder (rehost)	56	1	0	0	56	1	0
FaultFinder	30	88	110	1	6	1	0
Chaos Duck	10	47	0	0	0	0	57
ARCHIE	59	0	0	0	41	0	18
ARMORY	11	133	68	1	1	3	71

- Initialization state / Invalid access: inaccurate memory/register init. and memory mapping
 - → Rehosting
- Write attribute: Rehosting increased FPs due to the ability to correctly recognize memory regions
 - → Implementation of memory attributes
- Unimplemented instruction: ISE omits some specific instructions (e.g., ARM NEON).
 - → Manual verification

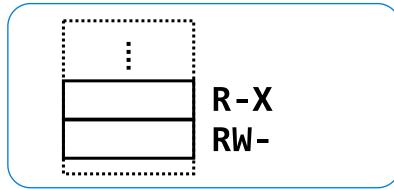
Key Takeaways



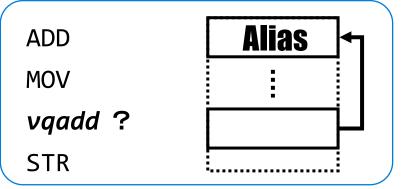
6 design principles for FVD



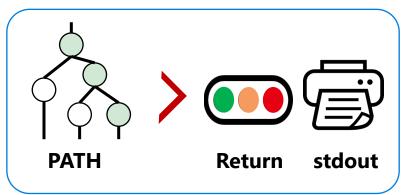
1. Initial state



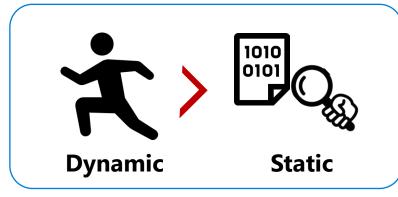
2. Memory map& attribute



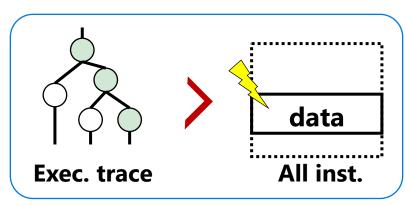
3. Unimplemented inst.& arch.-specific feature



4. Address reachability



5. Dynamic analysis



6. Execution trace

(success condition)

Applicability & Limitation



Ground truth depends on 1) 7 FVD tools capabilities and 2) execution platform

Rehosting Applicability

- Applicable to any ISE-based tools that support pre-execution memory modification
 - → ARMORY does not have such interface

Limitation

- 1. Tool dependency
 - Our ground truth is limited by the detection capabilities of the selected tools
- 2. Execution platform dependency
 - Our ground truth depends on I) native execution for x86, II) QEMU for ARM



Conclusion

Conclusion



Challenge: 1) Speed vs. Fidelity for FVD tools. 2) No benchmark dataset.

Contributions:

- Rehosting improved fidelity in ISE-based tool
- We created dataset "FIVBinBench" and it is available on GitHub[*]
- Comparative analysis of existing OSS-based FVD tools

Conclusion:

FaultFinder with rehosting offers the best balance of speed and performance



Future work

[*] https://github.com/pyth0n14n/FIVBinBench

- Implement memory attribute on FaultFinder
- More complex programs such as cryptography and VerifyPIN with countermeasure



Fault Injection Vulnerability Binary Benchmark

Main components

- Dataset including 1) binary and 2) result
- Tools including FaultFinder with rehosting

Instruction information

IP / code / opcode

Detection results

in each tool

Ground truth

with misclass. type

IP	code	opcode	oprand	comment	FaultARM	fault-injection-simulation	ChaosDuck	FaultFinder (rehosting)	FaultFinder	Grand Truth	Туре
00000598 <bytearraycompare>:</bytearraycompare>											
598:	e52db004	push	{fp}	@ (str fp, [sp, #-4]!)							
5b0:	e54b3015	strb	r3, [fp, #-21]			О	0	o	0	0	-
5b4:	e3a03000	mov	r3, #0		o	0	0	o	0	0	-
5b8:	e50b3008	str	r3, [fp, #-8]		o		0	o		0	Initial state
	00000598 <bytearraycompare>: 598: 560:</bytearraycompare>	00000598 598: e52db004 650: e54b3015 654: e3a03000	00000598 598: e52db004 push 5b0: e54b3015 strb 6b4: e3a03000 mov	00000598 698: e52db004 push {fp} 6b0: e54b3015 strb r3, [fp, #-21] 6b4: e3a03000 mov r3, #0	00000598 698: e52db004 push {fp} @ (str fp, [sp, #-4]!) 6b0: e54b3015 strb r3, [fp, #-21] 6b4: e3a03000 mov r3, #0	00000598 698: e52db004 push {fp} @ (str fp, [sp, #-4]!) 6b0: e54b3015 strb r3, [fp, #-21] o	00000598 698: e52db004 push {fp} @ (str fp, [sp, #-4]!) o 6b0: e54b3015 strb r3, [fp, #-21] o 6b4: e3a03000 mov r3, #0 o o	00000598 598: e52db004 push {fp} @ (str fp, [sp, #-4]!) o o o o o	00000598 598: e52db004 push {fp} @ (str fp, [sp, #-4]!)	00000598 000000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 000000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 000000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 00000598 000000598 00	00000598 598: e52db004 push {fp} @ (str fp, [sp, #-4]!) 5b0: e54b3015 strb r3, [fp, #-21] o o o o o o o o o o o o o o o o o o o

